



A second life for Prolog

Declarative programming

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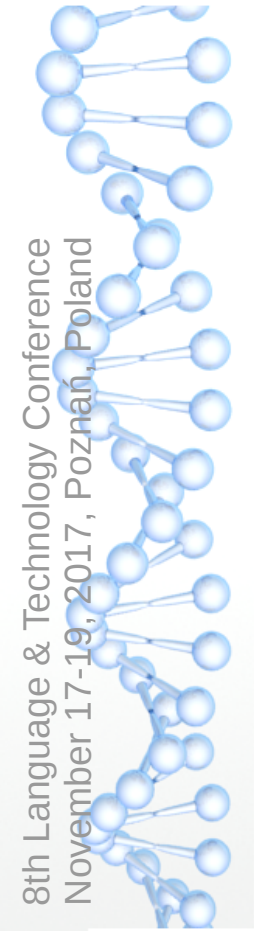
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Sessions

- 1) Introducing declarative programming and Prolog
 - Programming paradigms
 - Basics of Prolog
- 2) Algorithm = Logic + Control
 - Advanced control: tabling, constraints, continuations
 - Data and aggregation
- 3) Prolog as unifying framework
 - Accessing the outside world



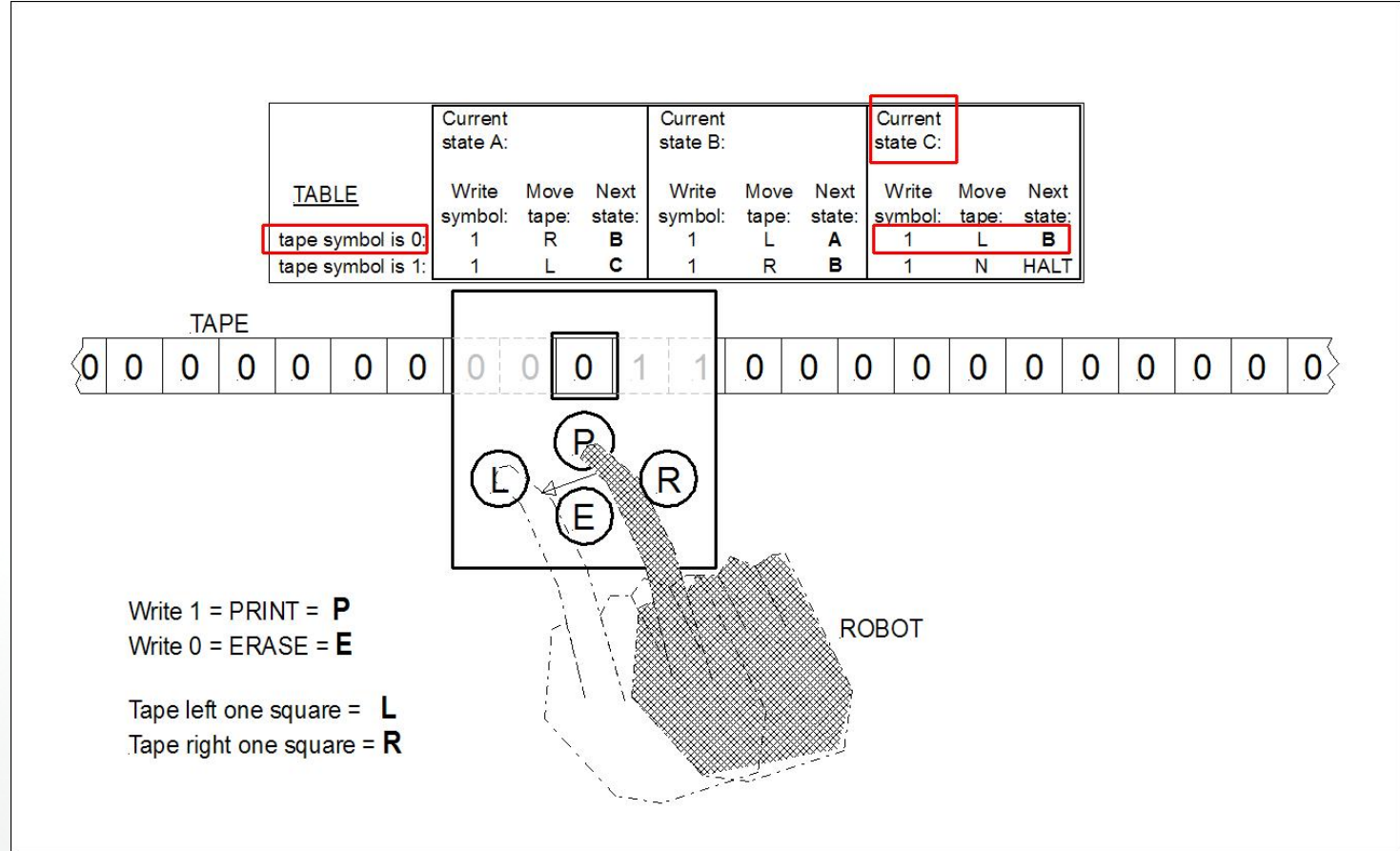


Programming paradigms





The Turing Machine



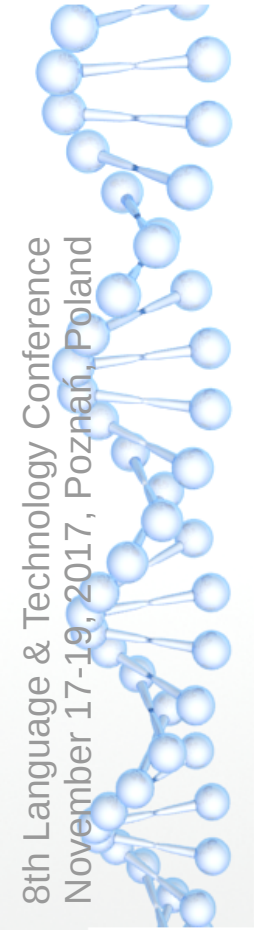
8th Language & Technology Conference
November 17-19, 2017, Poznan, Poland



The universal computing device!

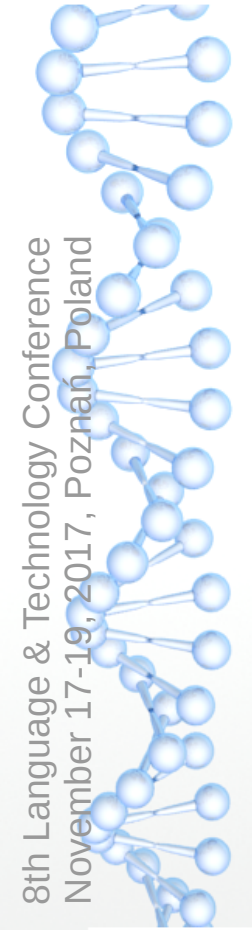
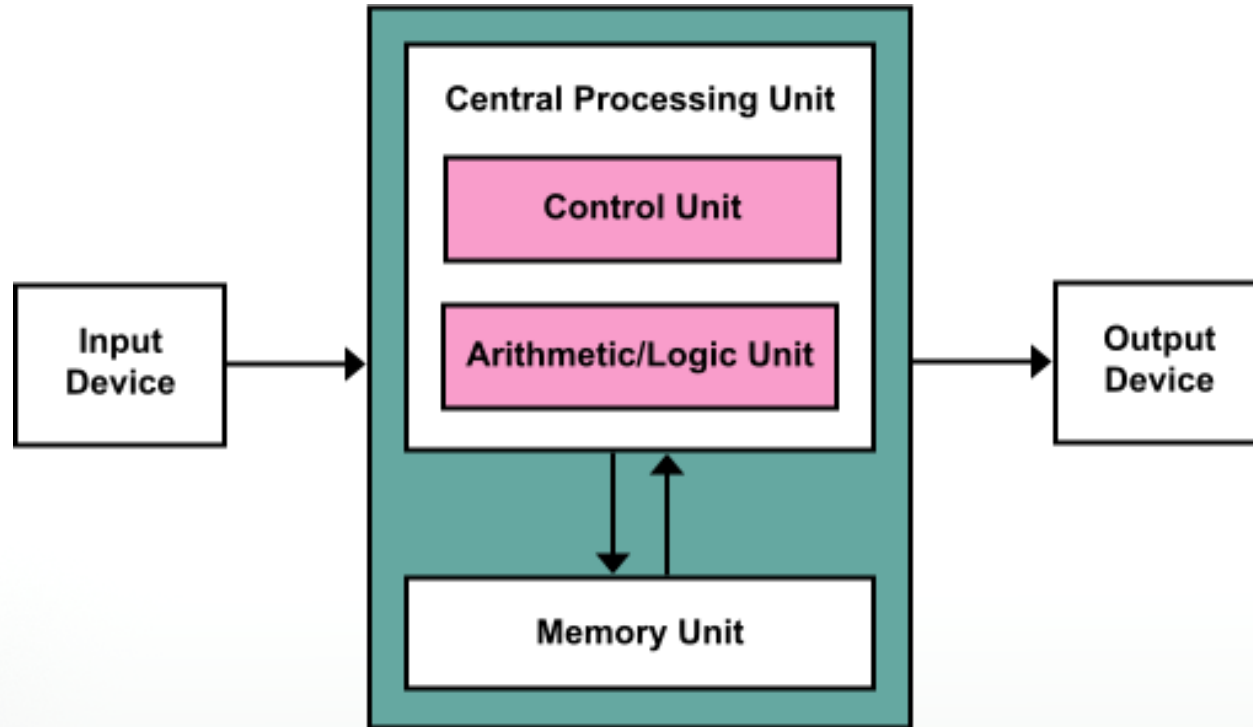
- "The evidence seems to indicate that every algorithm for any computing device has an equivalent Turing machine algorithm ... if [Church's thesis] is true, it is certainly remarkable that Turing machines, with their extremely primitive operations, are capable of performing any computation that any other device can perform, regardless of how complex a device we choose."

(Stone (1972), p. 13)





Von Neumann architecture





Imperative Programming

- Imperative: *giving an authoritative command*
- $X = Y + Z$
 - *Take the values stored in the memory locations of Y and Z, add them and store the result in the memory location of X.*
- $X = X + 1$
 - *To a **mathematician** this simply false*
 - *A **programmer** doesn't even see the problem!*
- *With state everywhere, it gets hard to*
 - *Understand the computation*
 - *Reorder it computation (exploit concurrency)*



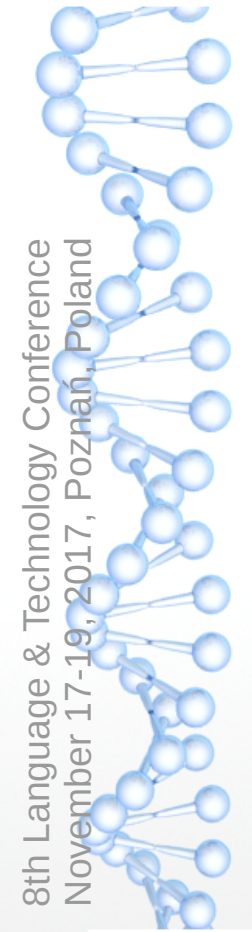


Just a test ...

```
do
{ double r = floor(e0/e1);
  double e00 = e0, p00 = p0, q00 = q0;
  volatile double p1_q1;

  e0 = e1;
  p0 = p1;
  q0 = q1;
  e1 = e00 - r*e1;
  p1 = p00 - r*p1;
  q1 = q00 - r*q1;

  p1_q1 = p1/q1;
  d = p1_q1 - n1->value.f;
} while(fabs(d) > DBL_EPSILON);
```





Functional programming

- A function specifies its (return) value as an expression over its inputs
 - There are no more variables storing state
 - Easier to reason about and execute concurrently
- Functions as primary objects (Lambda expressions)
- Lisp (1958), Scheme (1970), Clojure (2007)
- Most are hybrid language: they do provide for traditional variables but discourage their use
- JavaScript, Python. Even Java and to some extent C.



Logic programming

- Use (predicate) logic to describe the **relation** between values:

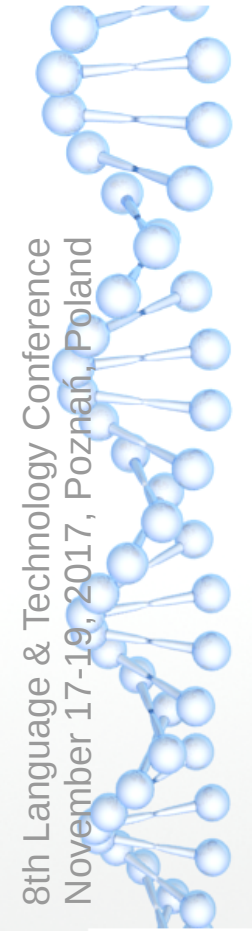
parent(bob, jane).

parent(jan, jane).

mother(X,Y) :- parent(X,Y), female(Y).

father(X,Y) :- parent(X,Y), male(Y).

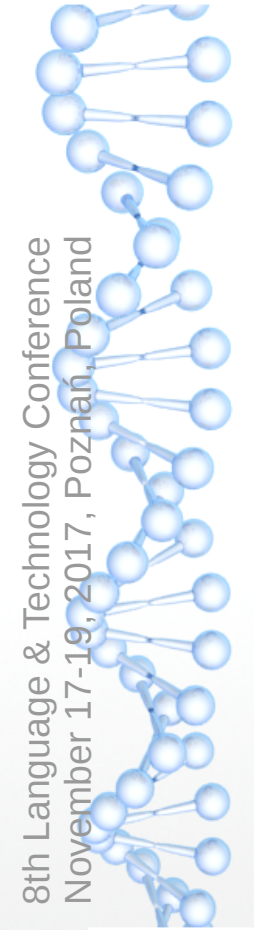
brother(X,Y) :- parent(X,P), parent(Y,P), male(Y).





Twists wrt. logic

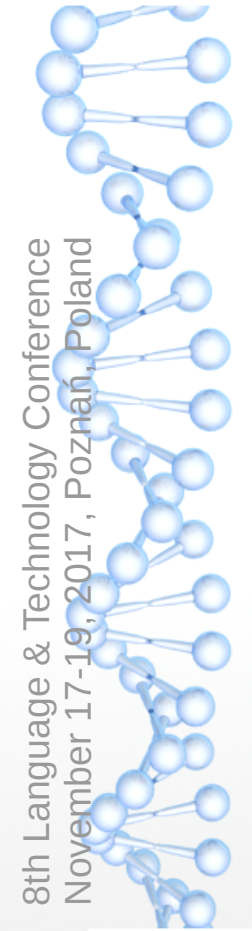
- $P \rightarrow Q$ (P implies Q) is written as $Q :- P$ (Q is true if P is true).
- Q (the head) is only a single *atom*, i.e., we cannot write
 - $A, B :- P.$





Some logic based languages

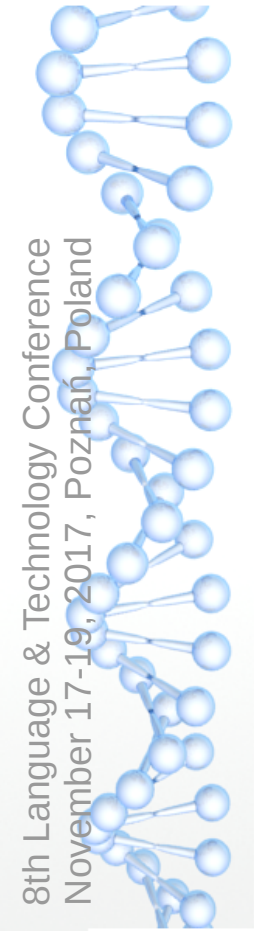
- Prolog (Alain Colmerauer, 1972)
 - Goal-directed backward chaining (top-down)
- Datalog (~1977)
 - Forward chaining (bottom-up)
- Mercury (1995)
 - Avoid need for extra-logical constructs using more declarations
- Picat (2015)
 - Aim at combinatorial problems
- Answer set programming (1993)
 - Pure





Prolog data

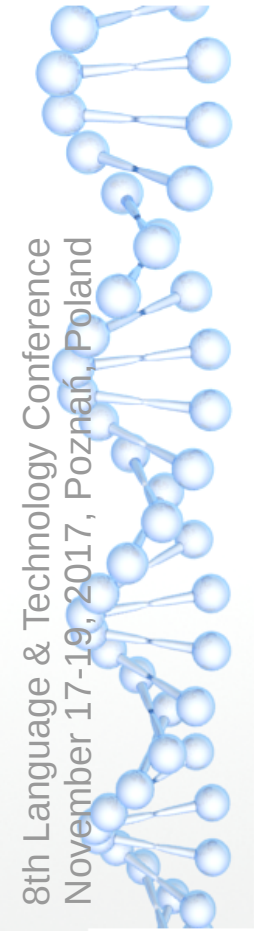
- Constants
 - bob, 'Bob', 'Nice wheather'
 - =, ==, <, ... (sequences of *glueing* punctuation characters)
 - !, |, ... (non-glueing punctuation characters)
- Numbers
 - 0, 42, 67.3
 - 7'012345, 0'a (97), 0b010101, 0xa0, 0o77
- Strings (SWI, LPA, ECLiPSe)
 - "abc" (SWI, ECLiPSe), `abc` (LPA)
- Compounds
 - age('Bob', 42)
- Variables
 - X, Myvar, _var





Syntactic sugar

- Operators (user extensible):
 - $X = Y \leftrightarrow =(X,Y)$
 - Type: **fx, fy, xfx, xfy, yfx, xf, yf**
 - Priority (0..1200)
- Lists (In SWI-Prolog `.` is `[]`):
 - $[a,b,c] \leftrightarrow .(a, .(b, .(c, [])))$
 - $[a,b|C] \leftrightarrow .(a, .(b, C))$
 - “ab“ or `ab` (SWI) $\leftrightarrow .(97,.(98, []))$
- Others
 - $\{ a,b \} \leftrightarrow \{ \{(a,b) \}$.





Logical variables: unification

- $T1 = T2$, is true if there is a substitution of variables with values that makes the two terms identical ($T1 == T2$)
- $X = a$
- $a = X$
- $X = Y$
- $X = Y, Y = a$
- $f(1) = f(X)$
- $1 = 2$
- $f(X,1) = f(Y,2)$
- $a = a$
- $a = a$
- $XY = XY$
- $a = a$
- $f(1) = f(1)$
- **false**
- **false**





Logical variables

- After unifying with a non-variable, the variable is *gone*!
 - ?- $T = f(X)$, $X = 1$.
 - $T = f(1)$.
- Unify-once:
 - $X = 1$, $X = 2 \rightarrow$ *false*!



Observations

- Objects (terms) that do not contain variables are **ground**.
- There is no **new**. Writing an object (term) creates it. As terms are **immutable one instance suffices**.
- The **program** is represented as a **term**
 - We can reason about programs
 - We can do dynamic programming



databases and a queries

- The data
 - mother(bob, jane).
 - mother(jan, jane).
- Queries
 - ?- mother(bob, jane).
 - **true.**
 - ?- mother(bob, jennifer).
 - **false.**



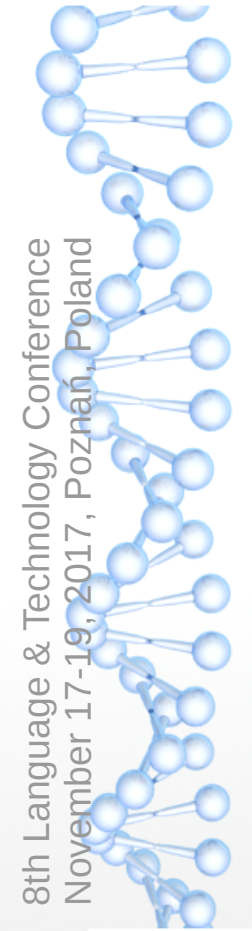
databases and a queries (cont)

- The data
 - mother(bob, jane).
 - mother(jan, jane).
- Queries
 - ?- mother(X, jane).
 - X = bob ;
 - X = jan ;
 - **false.**



databases and a queries (cont)

- The data
 - parent(jane, bob). female(jane). male(bob).
 - parent(jane, jan). male(jan).
 - parent(jane, mary). female(mary).
- Queries
 - ?- parent(jane, X), female(X).
 - X = mary.





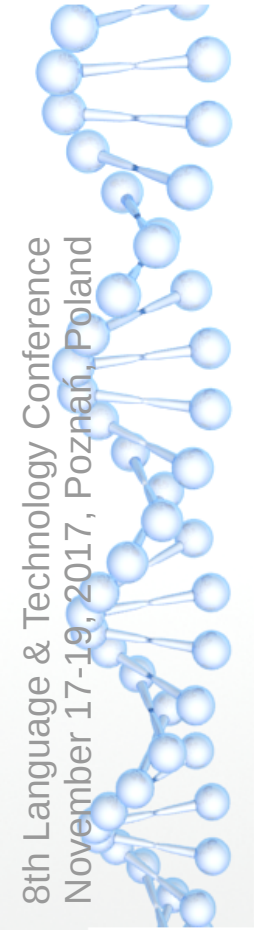
databases, rules and a queries

- The data
 - parent(jane, bob). female(jane). male(bob).
 - parent(jane, jan). male(jan).
 - parent(jane, mary). female(mary).
- Rule
 - daughter(Mother, Daughter) :-
 parent(Mother, Daughter),
 female(Daughter).
- Queries
 - ?- daughter(X, jane).
 - X = mary.



Exercises

- Please play around with the exercise below:
 - https://swish.swi-prolog.org/p/ltc_family.swinb
 -
- All material is linked from
 - <https://swish.swi-prolog.org/p/LTC2017.swinb>





Einstein's Riddle

- Also known as „*The zebra puzzle*“

- https://swish.swi-prolog.org/example/houses_puzzle.pl

1. Five colored houses in a row, each with an owner, a pet, cigarettes, and a drink.
2. The English lives in the red house.
3. The Spanish has a dog.
4. They drink coffee in the green house.
5. The Ukrainian drinks tea.
6. The green house is next to the white house.
7. The Winston smoker has a serpent.
8. In the yellow house they smoke Kool.
9. In the middle house they drink milk.
10. The Norwegian lives in the first house from the left.
11. The Chesterfield smoker lives near the man with the fox.
12. In the house near the house with the horse they smoke Kool.
13. The Lucky Strike smoker drinks juice.
14. The Japanese smokes Kent.
15. The Norwegian lives near the blue house.



**Five colored houses in a row,
each with an **owner**, a **pet**, **cigarettes**, and a
drink.**

```
Houses = [ house(Owner,Pet,Cigarettes,Drink,Color),  
            house(_____,_____,_____,_____,_____),  
            house(_____,_____,_____,_____,_____),  
            house(_____,_____,_____,_____,_____),  
            house(_____,_____,_____,_____,_____) ].
```

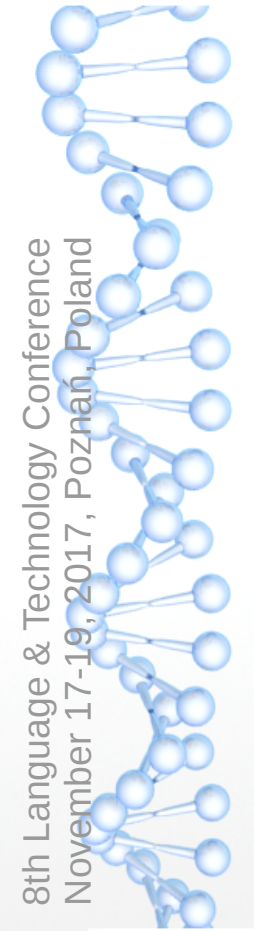




The English lives in the red house

```
member(house(english,_,_,_,red), Houses)
```

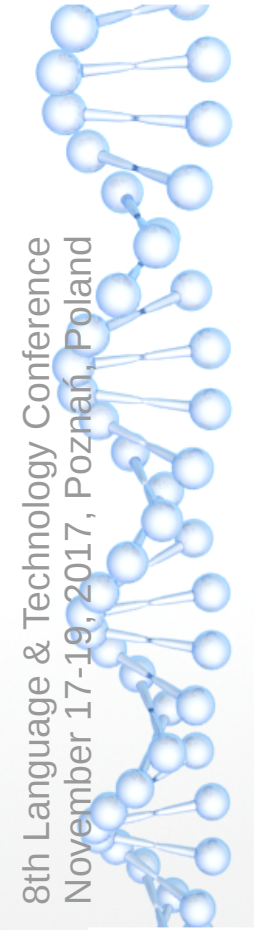
```
member(X, [X|_]).  
member(X, [_|T]) :- member(X, T).
```





In the middle house they drink milk.

Houses = [_,_,house(?,?,_,milk,_),_,_].

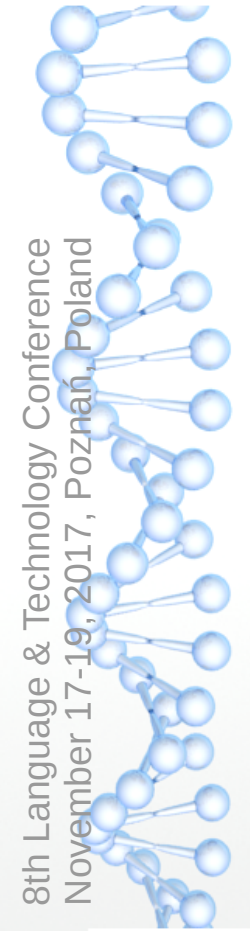




The Chesterfield smoker lives near the man with the fox

```
next(house(_,fox,_,_,_),  
     house(_,_,chesterfield,_,_),  
     Houses)
```

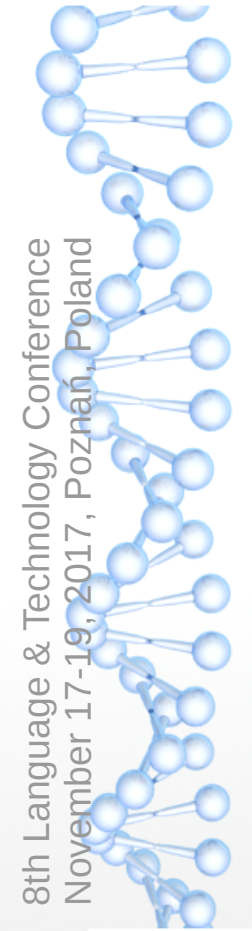
```
next(A, B, Ls) :- append(_, [A,B|_], Ls).  
next(A, B, Ls) :- append(_, [B,A|_], Ls).
```





Exercises

- Map coloring
 - https://swish.swi-prolog.org/p/ltc_mapcolor.swinb
- Learn Prolog Now! (2.2 Proof Search)
 - <http://lpn.swi-prolog.org/lpnpage.php?pagetype=html&pageid=lpn-htmlse6>





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